

# Fast and smooth simulation of space-time problems

## Day 3'



**Fast and smooth  
simulation of  
space-time problems**

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**Maciej Paszynski**  
Department of Computer Science, AGH University  
of Science and Technology, Krakow, Poland

Desde 24 al 28 de Julio, 2017  
Todos los días de 15:00 a 17:00 hrs.  
Sala Aula, Instituto de Matemáticas PUCV

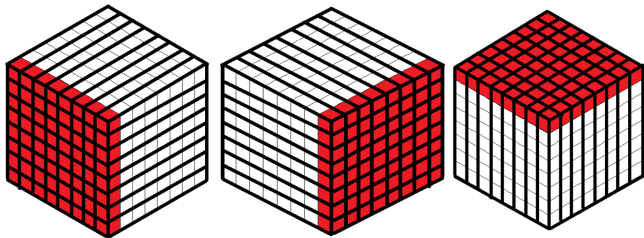
Department of Computer Science  
AGH University of Science and Technology, Kraków, Poland  
*[home.agh.edu.pl/paszynsk](http://home.agh.edu.pl/paszynsk)*

- Isogeometric finite element method
- Alternating Directions Implicit (ADI) method
- Isogeometric L2 projections
- Explicit dynamics
- Example 1: Heat transfer
- Installation of IGA-ADS solver
- Parallel shared memory explicit dynamics
- Example 2: Non-linear flow in heterogenous media
- **Parallel distributed memory explicit dynamics**
- Example 3: Linear elasticity
- Implicit dynamics
- Example 4: Implicit heat transfer
- Example 5: Pollution problem
- Labs with implicit dynamics

# Parallel domain decomposition based explicit dynamics

Parallel version for distributed-memory machines Message Passing Interface (**MPI**)

Maciej Woźniak, Marcin Łoś, Maciej Paszyński, Lisandro Dalcin, Victor Calo, *Parallel Fast Isogeometric Solvers for Explicit Dynamics*, Computing and Informatics 36(2) (2017)



**Figure:** Cube of processors. Gathers and scatters RHSs in three faces to perform three 1D solves with multiple RHS.

## Message Passing Interface (MPI)

subroutine **Gather**

(F, F\_out, n, elems, stride, dims, shifts, comm, ierr)

call **Linearize**(F\_out\_lin, F\_out, n+1, stride)

call **mpi\_gatherv** (F\_lin,    ← data to send by each processor 1, ...,  $N_x$   
elems \* stride,        ← size of data assigned to this processor  
MPI\_DOUBLE\_PRECISION,   ← type of data  
F\_out\_lin,              ← data to receive by processor 0 in a row  
dims,                   ← size of data from entire row  
shifts,                 ← offsets of slices from all processor from a row  
MPI\_DOUBLE\_PRECISION,   ← type of data to receive  
0, COMM, ierr)         ← communicator along a row

call **Delinearize**(F\_out\_lin, F\_out, n+1, stride)

# Parallel domain decomposition based explicit dynamics

Parallel version for distributed-memory machines (Fortran, MPI)

Maciej Woźniak, Marcin Łoś, Maciej Paszyński, Lisandro Dalcin, Victor Calo, *Parallel Fast Isogeometric Solvers for Explicit Dynamics*, Computing and Informatics 36(2) (2017)

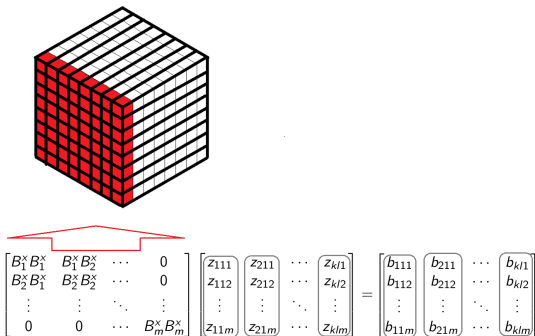


Figure: Gather and scatter data into three faces on cube of processors

Parallel version for distributed-memory machines (Fortran, MPI)

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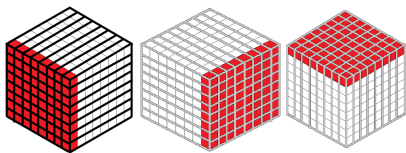


Figure: Gather into OYZ face

0. Integration

1a. Gather in every row of processors into OYZ face

1b. Solve  $N_y N_z$  1D problems with multiple RHS

1c. Scatter results onto cube of processors

1d. Reorder right hand sides

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Parallel version for distributed-memory machines (Fortran, MPI)

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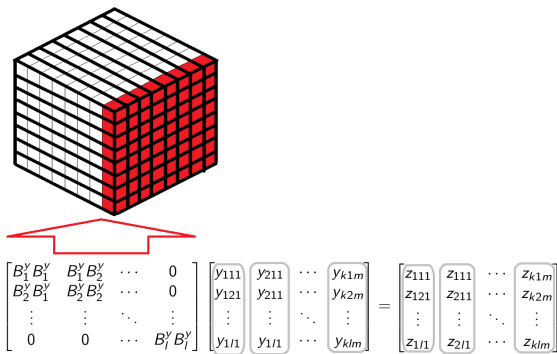


Figure: Gather into OXZ face

Parallel version for distributed-memory machines (Fortran, MPI)

Maciej Woźniak, Marcin Łoś, Maciej Paszyński, Lisandro Dalcin, Victor Calo, *Parallel Fast Isogeometric Solvers for Explicit Dynamics*, Computing and Informatics 36(2) (2017)

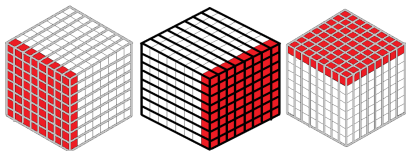


Figure: Gather and scatter data into three faces on cube of processors

- 2a. Gather in every row of processors into  $OXZ$  face
- 2b. Solve  $N_x N_z$  1D problem with multiple RHS
- 2c. Scatter results onto cube of processors
- 2d. Reorder right hand sides



# Parallel domain decomposition based explicit dynamics

Parallel version for distributed-memory machines (Fortran, MPI)

Maciej Woźniak, Marcin Łoś, Maciej Paszyński, Lisandro Dalcin, Victor Calo, *Parallel Fast Isogeometric Solvers for Explicit Dynamics*, Computing and Informatics 36(2) (2017)

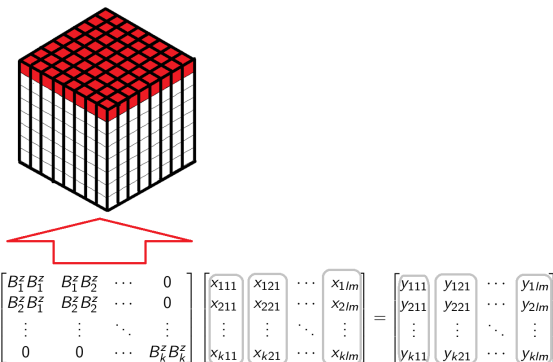


Figure: Gather into OXY face

## Parallel version for distributed-memory machines (Fortran, MPI)

Maciej Woźniak, Marcin Łoś, Maciej Paszyński, Lisandro Dalcin, Victor Calo, *Parallel Fast Isogeometric Solvers for Explicit Dynamics*, Computing and Informatics 36(2) (2017)

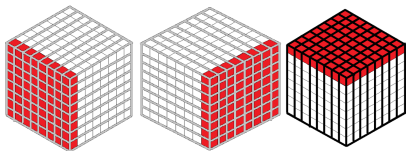


Figure: Gather and scatter data into three faces on cube of processors

- 3a. Gather in every row of processors into  $OXY$  face
- 3b. Solve  $N_x N_y$  1D problem with multiple RHS
- 3c. Scatter results onto cube of processors
- 3d. Reorder right hand sides

# Parallel domain decomposition based explicit dynamics

We have a mesh of  $N_x \times N_y \times N_z$  elements

There are  $(p_x + 1)(p_y + 1)(p_z + 1)$  basis functions over each element

Integration over all elements  $O\left(\frac{p_x^2 p_y^2 p_z^2 N_x N_y N_z}{c_x c_y c_z}\right)$

Solution (1D problem multiple RHS)  $O\left(\frac{(p_x^2 c_x + p_y^2 c_y + p_z^2 c_z)(N_x N_y N_z)}{c_x c_y c_z}\right)$

Gather  $O\left(\frac{(c_x + c_y + c_z) N_x N_y N_z}{c_x c_y c_z}\right)$

Reorder  $O\left(\frac{N_x N_y N_z p_x p_y p_z}{c_x c_y c_z}\right)$

Scather  $O\left(\frac{(c_x + c_y + c_z) N_x N_y N_z}{c_x c_y c_z}\right)$

Assuming

$$N_x = N_y = N_z = N^{1/3}, \quad p_x = p_y = p_z = p, \quad c_x = c_y = c_z = c^{1/3}$$

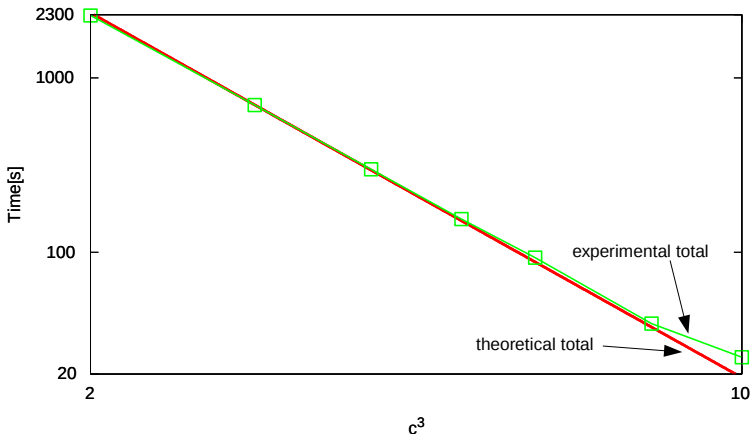
we have the following cost  $\left(\frac{p^6 N}{c} + \frac{p^2 N}{c^{2/3}} + \frac{p^3 N}{c}\right) t_{comp} + \left(\frac{N}{c^{2/3}}\right) t_{comm}$

which implies **the computational complexity**  $O\left(p^6 \frac{N}{c}\right)$

and **the communication complexity**  $O\left(\frac{N}{c^{2/3}}\right)$

# Parallel domain decomposition based explicit dynamics

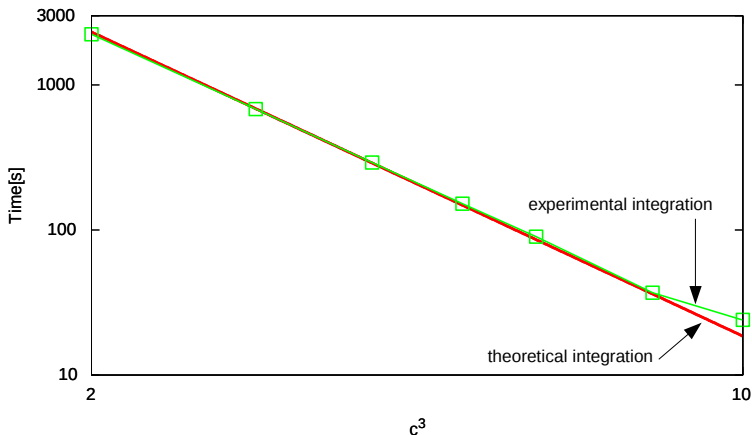
## Numerical experiments LONESTAR Linux cluster



**Figure:** Comparison of total experimental and theoretical execution time for  $N = 512$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

The integration time is dominating the solution time significantly.



**Figure:** Comparison of experimental and theoretical integration time for  $N = 512$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

The solution time takes around 1 percent of the total solver time.

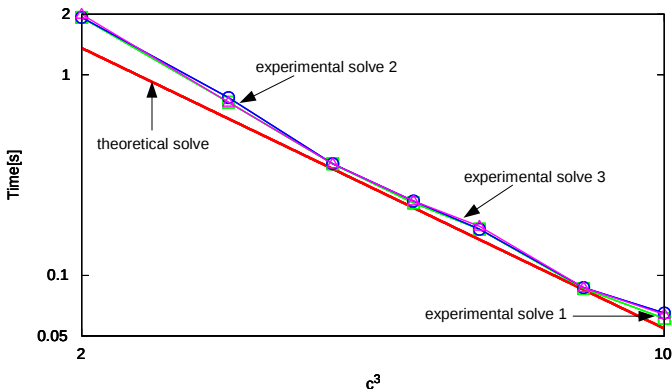


Figure: Comparison of experimental and theoretical solution times for  $N = 512$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

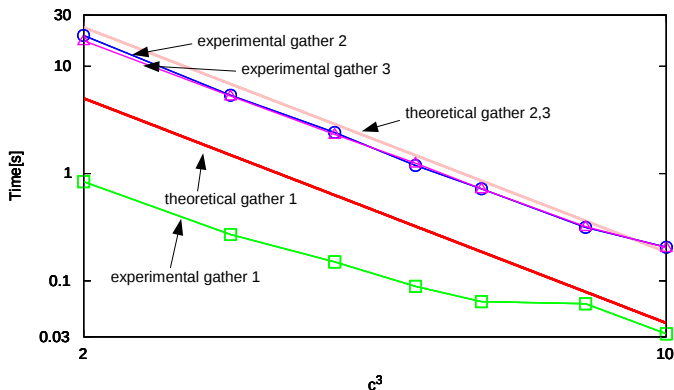
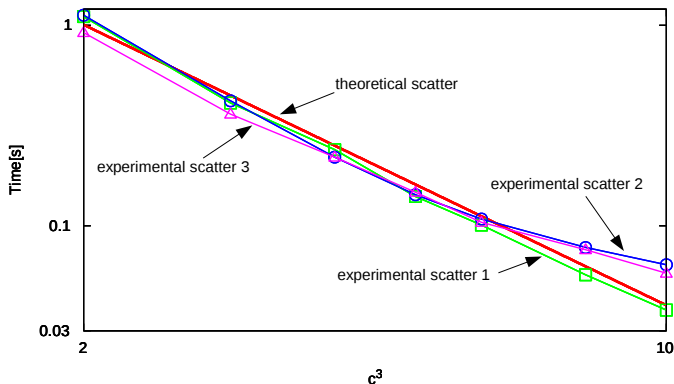


Figure: Comparison of experimental and theoretical gather times for  $N = 512$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

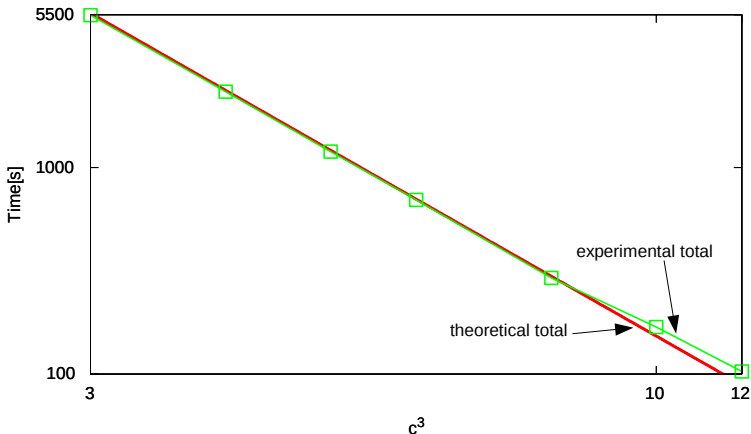


**Figure:** Comparison of experimental and theoretical scatter times for  $N = 512$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .



# Parallel domain decomposition based explicit dynamics

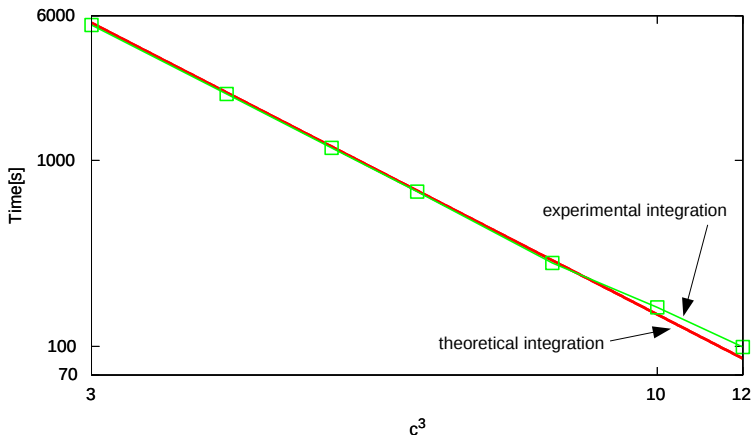
## Numerical experiments LONESTAR Linux cluster



**Figure:** Comparison of total experimental and theoretical execution time for  $N = 1024$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

The integration time is dominating the solution time significantly.



**Figure:** Comparison of total experimental and estimated integration time for  $N = 1024$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

The solution time takes around 1 percent of the total solver time.

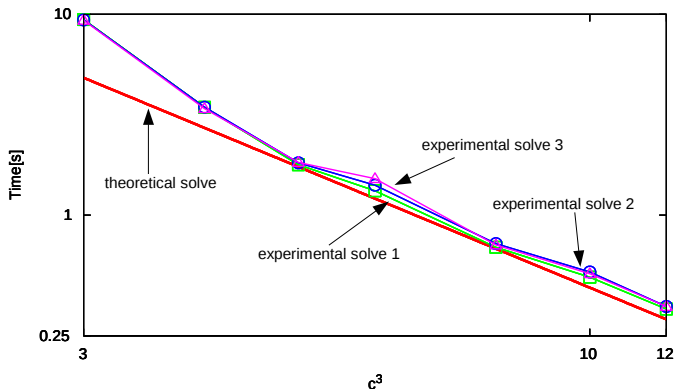


Figure: Comparison of total experimental and estimated solution times for  $N = 1024$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

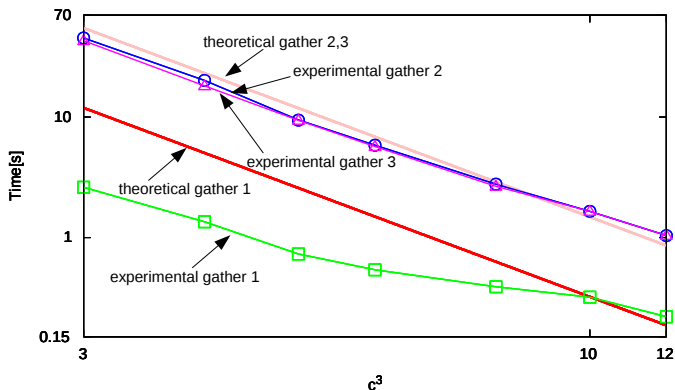


Figure: Comparison of total experimental and estimated gather times for  $N = 1024$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Parallel domain decomposition based explicit dynamics

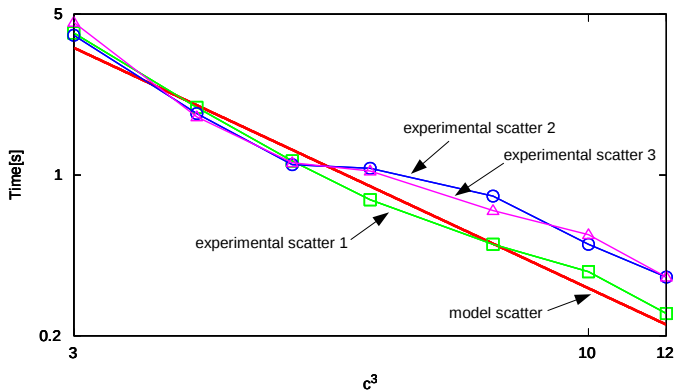


Figure: Comparison of total experimental and estimated scatter times for  $N = 1024$  for  $p = 3$  for different number of processors  $c^3 = 2^3, \dots, 10^3 = 8, \dots, 1000$ .

# Limitations of isogeometric L2 projections

Time step size limited by Courant-Fredrichs-Lewy (CFL) condition

[https://en.wikipedia.org/wiki/Courant-Friedrichs-Lewy\\_condition](https://en.wikipedia.org/wiki/Courant-Friedrichs-Lewy_condition)

$$\frac{u_x \Delta t}{\Delta x} + \frac{u_y \Delta t}{\Delta y} \leq C_{max}$$

where  $\Delta t$  is time step size,  $\Delta x = \Delta y = h$  element size,

$u_x, u_y$  magnitude of the field,  $C_{max} = 1$  for explicit method, so

$$(u_x + u_y) * \Delta t \leq h$$

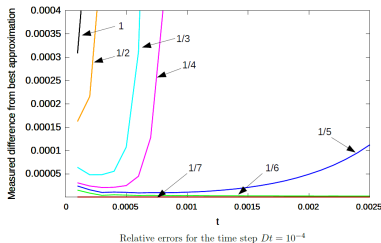


Figure: Lack of convergence for  $Dt = 10^{-4}, \frac{10^{-4}}{2}, \dots, \frac{10^{-4}}{5}$

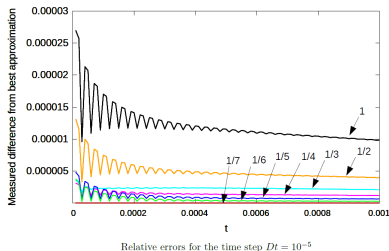


Figure: Convergence for  $Dt = 10^{-5}$  and smaller time steps

# Example: Too large time step for elastic wave propagation

Click in the middle

## Example: Good time step for elastic wave propagation

Click in the middle



# Integration of the 3D matrix

If we go for implicit method, we need to integrate the matrix

We have a mesh of  $N_x \times N_y \times N_z$  elements

$nrdof = (p_x + 1)(p_y + 1)(p_z + 1)$  basis functions on each element

$(p_x, p_y, p_z)$  denotes the B-splines order in directions  $x, y$  and  $z$

$ngx = O(p_x)$ ,  $ngy = O(p_y)$ ,  $ngz = O(p_z)$  number of Gauss points

$$\begin{aligned} & (B_{i,p}^x B_{j,p}^y B_{k,p}^z, B_{l,p}^x B_{m,p}^y B_{n,p}^z)_{L2} = \\ & \int_{\Omega} B_{i,p}^x(x) B_{j,p}^y(y) B_{k,p}^z(z) B_{l,p}^x(x) B_{m,p}^y(y) B_{n,p}^z(z) dx dy dz = \\ & \sum_E \int_E B_{i,p}^x(x) B_{j,p}^y(y) B_{k,p}^z(z) B_{l,p}^x(x) B_{m,p}^y(y) B_{n,p}^z(z) dx dy dz = \\ & \sum_E \sum_{s=1,ngx;t=1,ngy;w=1,ngz} W_s W_t W_z B_{i,p}^x(x_s) B_{j,p}^y(y_t) B_{k,p}^z(z_w) B_{l,p}^x(x_s) B_{m,p}^y(y_t) B_{n,p}^z(z_w) \end{aligned}$$

We construct the element matrices for each element, for all the basis functions which span over the element, namely

$B_{i,p}^x, i = 1, p_x + 1$ ,  $B_{j,p}^y, j = 1, p_y + 1$ ,  $B_{k,p}^z, k = 1, p_z + 1$  (trial functions) and  $B_{l,p}^x, l = 1, p_x + 1$ ,  $B_{m,p}^y, m = 1, p_y + 1$ ,  $B_{n,p}^z, n = 1, p_z + 1$  (test functions).

# Integration of the 3D matrix

We have a mesh of  $N_x \times N_y \times N_z$  elements

$nrdof = (p_x + 1)(p_y + 1)(p_z + 1)$  basis functions on each element

$(p_x, p_y, p_z)$  denotes the B-splines order in directions  $x, y$  and  $z$

$ngx = O(p_x), ngy = O(p_y), ngz = O(p_z)$  number of Gauss points

- for  $s = 1, ngx$  // Gauss integration points

- for  $t = 1, ngy$

- for  $w = 1, ngz$

get Gauss point  $(x_s, y_t, z_w)$ , weight  $W_s W_t W_w$

- for  $l = 1, p_x + 1$  // trial B-splines

- for  $m = 1, p_y + 1$

- for  $n = 1, p_z + 1$

- for  $j = 1, p_x + 1$  // test B-splines

- for  $j = 1, p_y + 1$

- for  $k = 1, p_z + 1$

- aggregate

$$W * B_{i,p}^x(x_s) B_{j,p}^y(x_t) B_{k,p}^z(x_w), B_{l,p}^x(x_s) B_{m,p}^y(x_t) B_{n,p}^z(x_w)$$

$p_x = p_y = p_z = p$  computational complexity  $O(p^9)$ , if  $p=9$  it is  $10^9$

# Sequential integration of the 3D matrix

$F = 0.d0$

```
do ex = 1,nelemx //Loop through elements
  do ey = 1,nelemy
    do ez = 1,nelemz
      J = Jx(ex)*Jy(ey)*Jz(ez) //element Jacobian
      do kx = 1,ngx //Loop through Gauss points
        do ky = 1,ngy
          do kz = 1,ngz
            W = Wx(kx)*Wy(ky)*Wz(kz) //Gauss weight
            value = fvalue(Xx(kx,ex),Xy(ky,ey),Xz(kz,ez))
            do ax = 0,px //B-splines
              do ay = 0,py
                do az = 0,pz
                  do bx = 0,px //B-splines
                    do by = 0,py
                      do bz = 0,pz
                        call compute_index(ind,ax,ay,az,ex,ey,ez,nx,ny,nz)
                        call compute_index(ind1,bx,by,bz,ex,ey,ez,nx,ny,nz)
                        A(ind,ind1) = A(ind,ind1) +
NNx(0,ax,kx,ex)*NNy(0,ay,ky,ey)*NNz(0,az,kz ,ez)*
NNx(0,bx,kx,ex)*NNy(0,by,ky,ey)*NNz(0,bz,kz ,ez)*J*W*value
```

# Parallel OpenMP integration of the 3D matrix

## OpenMP = Open Multi-Processing

```
!$OMP PARALLEL DO
!$OMP& DEFAULT(SHARED)
!$OMP& FIRSTPRIVATE
    (iy,ex,ey,ez,J,kx,ky,kz,W,value,ax,ay,az,bx,by,bz,ind,ind1)
!$OMP& REDUCTION(+:nr_nonzeros)
do iy=1,miy //Now it is 1 loop over elements
call map_indexes(iy,ex,ey,ez)
J = Jx(ex)*Jy(ey)*Jz(ez) //element Jacobian
do kx = 0,ngx //loop through Gauss points
do ky = 0,ngy
do kz = 0,ngz
W = Wx(kx)*Wy(ky)*Wz(kz) //Gauss weight
value = fvalue(Xx(kx,ex),Xy(ky,ey),Xz(kz,ez))
do ax = 0,px //trail B-splines along x,y,z
do ay = 0,py
do az = 0,pz
do bx = 0,px //test B-splines along x,y,z
do by = 0,py
do bz = 0,pz
call compute_index(ind,ax,ay,az,ex,ey,ez,nx,ny,nz)
call compute_index(ind1,bx,by,bz,ex,ey,ez,nx,ny,nz)
A(ind,ind1) = A(ind,ind1) +
NNx(0,ax,kx,ex)*NNy(0,ay,ky,ey)*NNz(0,az,kz,ez)*
NNx(0,bx,kx,ex)*NNy(0,by,ky,ey)*NNz(0,bz,kz,ez)*J*W*value
!$OMP END PARALLEL DO
```

# Numerical results

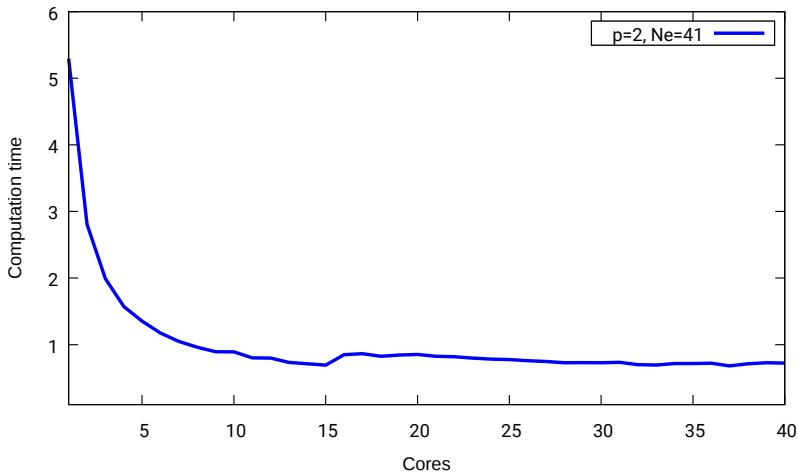


Figure: Execution time of the parallel integration algorithm, when increasing number of cores. 3D element with quadratic polynomials

# Numerical results

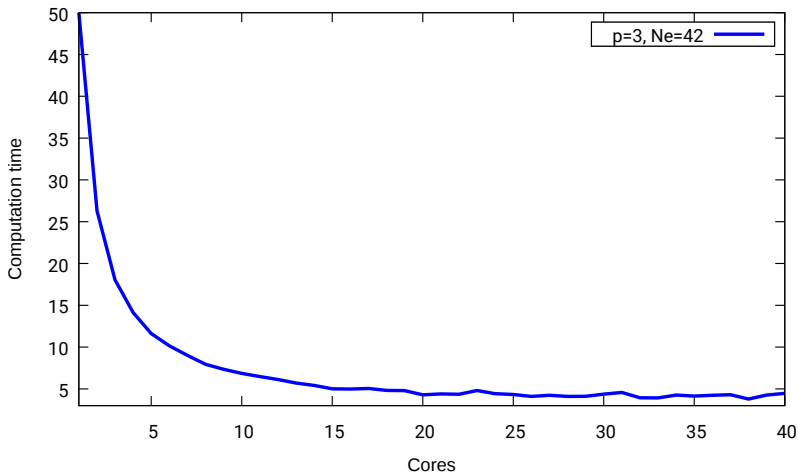


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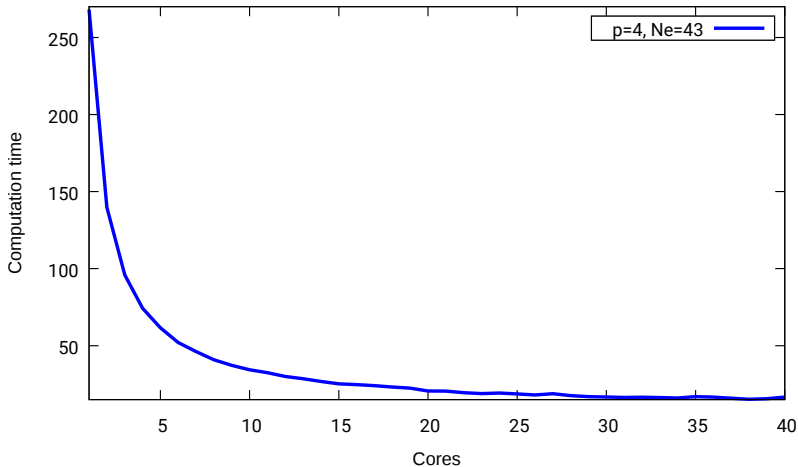


Figure: Execution time of the parallel integration algorithm, when increasing number of cores. 3D element with quartic polynomials

# Numerical results

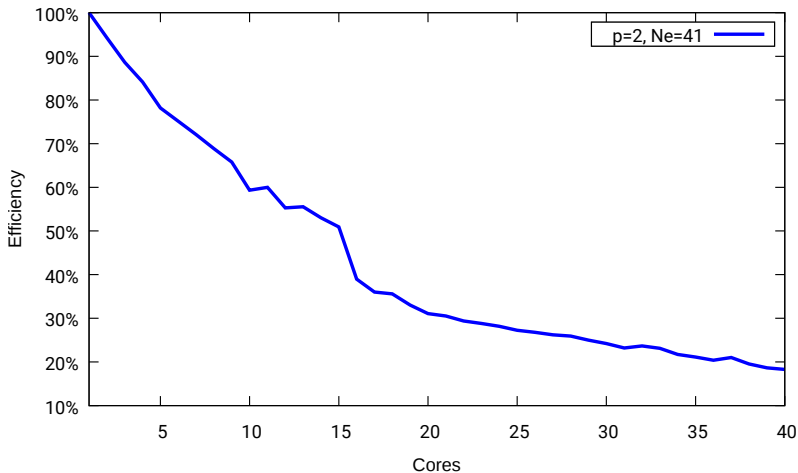


Figure: Efficiency of the parallel integration algorithm, when increasing number of cores. 3D element with quadratic polynomials



# Numerical results

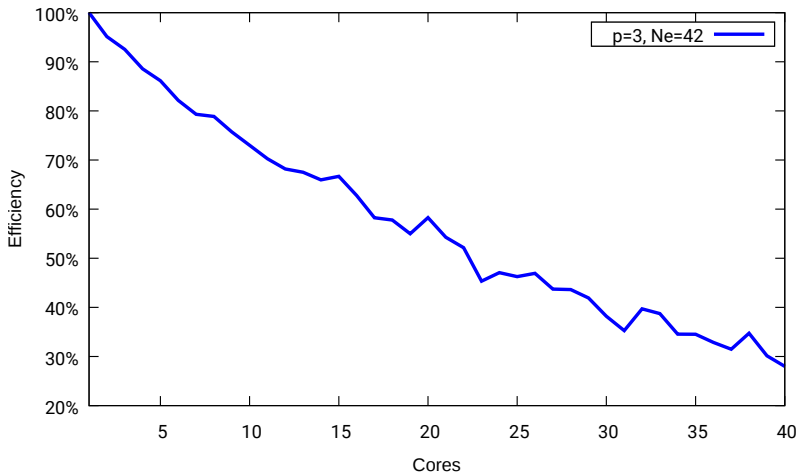


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# Numerical results

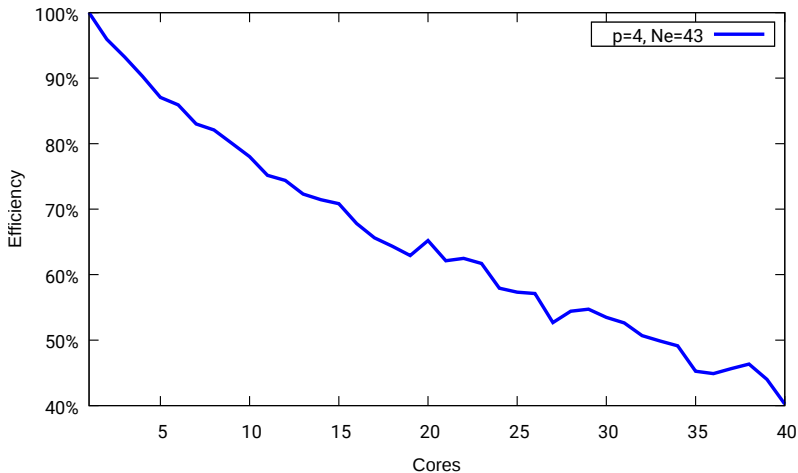


Figure: Efficiency of the parallel integration algorithm, when increasing number of cores. 3D element with quartic polynomials

# Numerical results

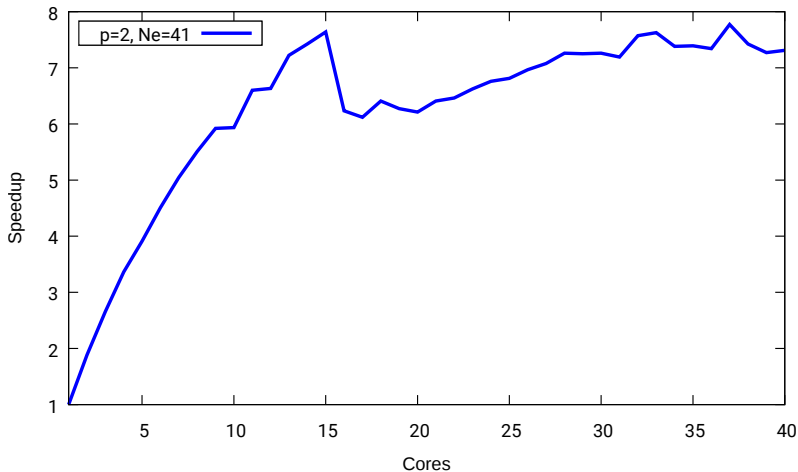


Figure: Speedup of the parallel integration algorithm, when increasing number of cores. 3D element with quadratic polynomials

# Numerical results

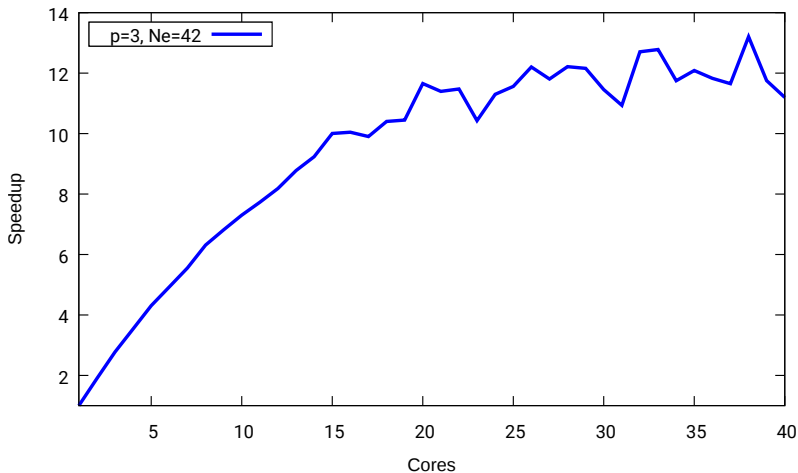


Figure: Speedup of the parallel integration algorithm, when increasing number of cores. 3D element with cubic polynomials

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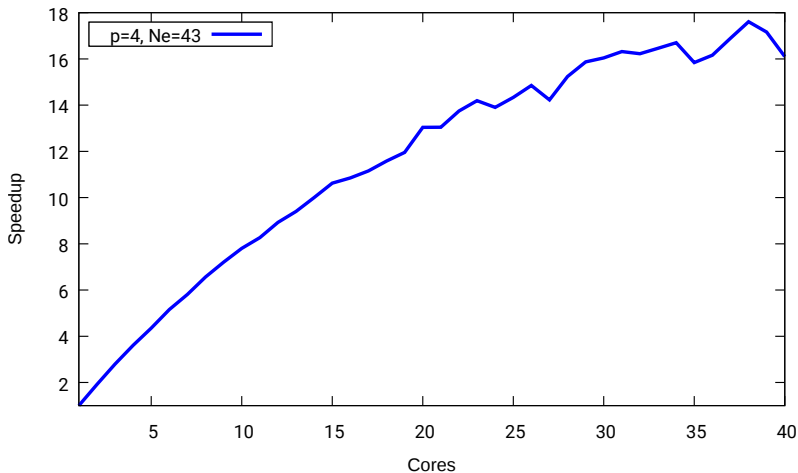


Figure: Speedup of the parallel integration algorithm, when increasing number of cores. 3D element with quartic polynomials